

Machine Prog: Data & Advanced

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Outline

- Array
- Structure
- Floating Point
- *Additional Exercise : Pointers & Dereference
- 可变长栈帧
- Buffer Overflow
- *Memory Layout

定长数组

循环遍历定长数组时，编译器会进行优化！

```
/* Compute i,k of fixed matrix product */
int fix_prod_ele (fix_matrix A, fix_matrix B, long i, long k) {
    long j;
    int result = 0;

    for (j = 0; j < N; j++)
        result += A[i][j] * B[j][k];

    return result;
}
```

```
#define N 16
typedef int fix_matrix[N][N];
```

原始代码

```
/* Compute i,k of fixed matrix product */
int fix_prod_ele_opt(fix_matrix A, fix_matrix B, long i, long k) {
    int *Aptr = &A[i][0];      /* Points to elements in row i of A */
    int *Bptr = &B[0][k];      /* Points to elements in column k of B */
    int *Bend = &B[N][k];     /* Marks stopping point for Bptr */
    int result = 0;
    do {                      /* No need for initial test */
        result += *Aptr * *Bptr; /* Add next product to sum */
        Aptr++;                 /* Move Aptr to next column */
        Bptr += N;               /* Move Bptr to next row */
    } while (Bptr != Bend);    /* Test for stopping point */
    return result;
}
```

优化后的代码

定长数组

访问定长的高维数组时，编译器会进行优化！

```
#define N 16  
typedef int fix_matrix[N][N];
```

```
/* Get element A[i][j] */  
int fix_ele(fix_matrix A, size_t i, size_t j) {  
    return A[i][j];  
}
```

```
# A in %rdi, i in %rsi, j in %rdx  
salq    $6, %rsi           # 64*i  
addq    %rsi, %rdi         # A + 64*i  
movl    (%rdi,%rdx,4), %eax # Mem[A + 64*i + 4*j]  
ret
```

可变长数组

```
int var_ele(long n, int A[n][n], long i, long j) {  
    return A[i][j];  
}
```

int var_ele(long n, int A[n][n], long i, long j)

n in %rdi, A in %rsi, i in %rdx, j in %rcx

1 var_ele:

2	imulq %rdx, %rdi	<i>Compute n · i</i>
3	leaq (%rsi,%rdi,4), %rax	<i>Compute $x_A + 4(n \cdot i)$</i>
4	movl (%rax,%rcx,4), %eax	<i>Read from M[$x_A + 4(n \cdot i) + 4j$]</i>
5	ret	

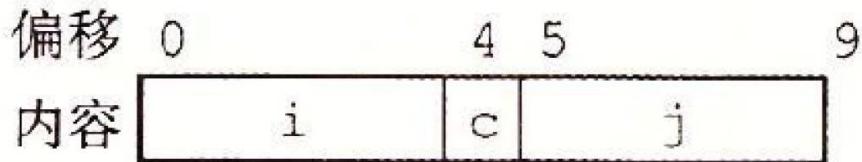
由于不定长，编译器不能用sal, leaq等指令加速乘法，被迫使用较慢的imulq

Structure

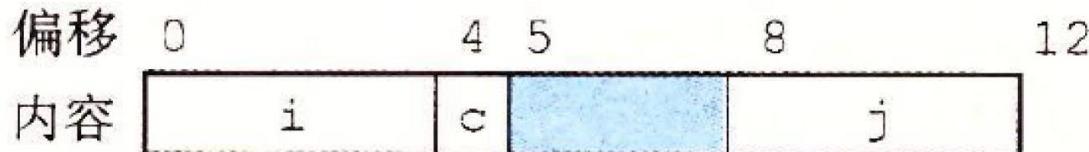
```
struct S1 {  
    int i;  
    char c;  
    int j;  
};
```

- 对于x86机器而言，处理未对齐的数据仍可正常运行，但是会影响效率；
- 对于其他的一些机器而言，处理未对齐的数据可能导致内存错误。

- 对齐前：



- 对齐后：



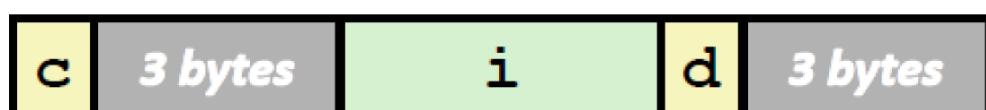
- ! 注意结构的末尾可能需要对齐（结构数组）
- ! 注意栈上的参数值总是8字节对齐

Structure

Saving Space: 贪心思想

```
struct S4 {  
    char c;  
    int i;  
    char d;  
} *p;
```

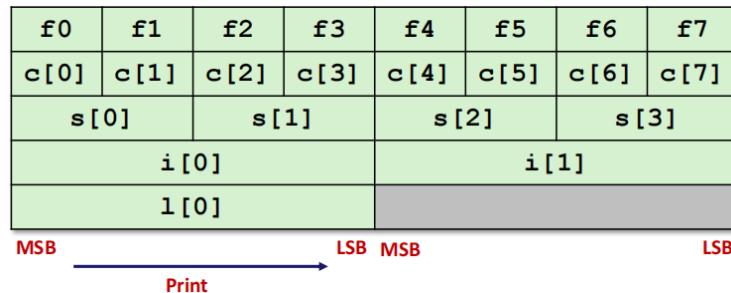
```
struct S5 {  
    int i;  
    char c;  
    char d;  
} *p;
```



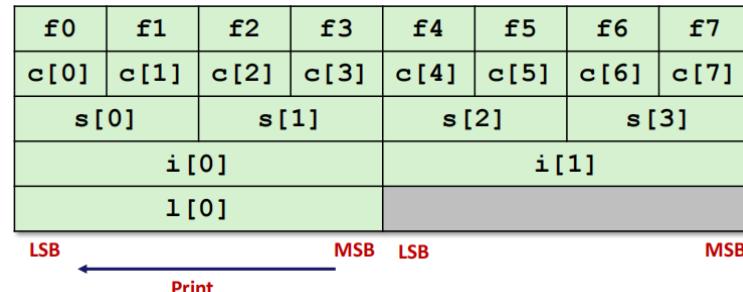
Union

Union中所有成员都从低地址开始存放，可据此判断大端/小端

Big Endian



Little Endian



Output:

```
Characters 0-7 == [0xf0,0xf1,0xf2,0xf3,0xf4,0xf5,0xf6,0xf7]
Shorts    0-3 == [0xf1f0,0xf3f2,0xf5f4,0xf7f6]
Ints     0-1 == [0xf3f2f1f0,0xf7f6f5f4]
Long      0 == [0xf3f2f1f0]
```

Output on Sun:

```
Characters 0-7 == [0xf0,0xf1,0xf2,0xf3,0xf4,0xf5,0xf6,0xf7]
Shorts    0-3 == [0xf0f1,0xf2f3,0xf4f5,0xf6f7]
Ints     0-1 == [0xf0f1f2f3,0xf4f5f6f7]
Long      0 == [0xf0f1f2f3]
```

Floating Point

%xmm0 返回值

%xmm0~7 8个参数

%xmm0~15 所有均为caller saved

这部分了解即可，
往年没有深入考察过

255	127	0
%ymm0	%xmm0	1st FP arg./Return value
%ymm1	%xmm1	2nd FP argument
%ymm2	%xmm2	3rd FP argument
%ymm3	%xmm3	4th FP argument
%ymm4	%xmm4	5th FP argument
%ymm5	%xmm5	6th FP argument
%ymm6	%xmm6	7th FP argument
%ymm7	%xmm7	8th FP argument
%ymm8	%xmm8	Caller saved
%ymm9	%xmm9	Caller saved
%ymm10	%xmm10	Caller saved
%ymm11	%xmm11	Caller saved
%ymm12	%xmm12	Caller saved
%ymm13	%xmm13	Caller saved
%ymm14	%xmm14	Caller saved
%ymm15	%xmm15	Caller saved

```
float fadd(float x, float y)
{
    return x + y;
}
```

```
double dadd(double x, double y)
{
    return x + y;
}
```

```
# x in %xmm0, y in %xmm1
addss    %xmm1, %xmm0
ret
```

```
# x in %xmm0, y in %xmm1
addsd    %xmm1, %xmm0
ret
```

Additional Exercise : Pointers & Dereference

	sizeof (A)	What is A?
int *A[3];		
int *(A[3]);		
int (*A)[3];		
int (*A[3]);		
int (*A[3])();		
int (*A[3])[5];		

Hint:

Operator	Read as	Priority
type A(parameters)	Function with parameters and return type type	0
type [N]	Array of N type	1
*	Pointer to	2

Additional Exercise : Pointers & Dereference

	<code>sizeof (A)</code>	What is A?
<code>int *A[3];</code>	24	Array of 3 int*
<code>int *(A[3]);</code>	24	Array of 3 int*
<code>int (*A)[3];</code>	8	Pointer to array of 3 int
<code>int (*A[3]);</code>	24	Array of 3 int*
<code>int (*A[3])();</code>	24	Array of 3 pointers to a function with no parameter and return type int
<code>int (*A[3])[5];</code>	24	Array of size 3 pointers to array of size 5 of int

Hint:

Operator	Read as	Priority
<code>type A(parameters)</code>	Function with parameters and return type type	0
<code>type [N]</code>	Array of N type	1
<code>*</code>	Pointer to	2

可变长栈帧

```
long vframe(long n, long idx, long *q)
n in %rdi, idx in %rsi, q in %rdx
Only portions of code shown
```

1 vframe:

```
2     pushq    %rbp
3     movq    %rsp, %rbp
4     subq    $16, %rsp
5     leaq    22(%rdi,8), %rax
6     andq    $-16, %rax
7     subq    %rax, %rsp
8     leaq    7(%rsp), %rax
9     shrq    $3, %rax
10    leaq    0(%rax,8), %r8
11    movq    %r8, %rcx
12
13    . . .
14
15    Code for initialization loop
16    i in %rax and on stack, n in %rdi, p in %rcx, q in %rdx
17    .L3:
18        movq    %rdx, (%rcx,%rax,8)
19        addq    $1, %rax
20        movq    %rax, -8(%rbp)
21
22    .L2:
23        movq    -8(%rbp), %rax
24        cmpq    %rdi, %rax
25        jl     .L3
26
27    . . .
28
29    Code for function exit
30
31    leave
32    ret
```

Save old %rbp
Set frame pointer }

Allocate space for i (%rsp = s₁)

Allocate space for array p (%rsp = s₂)

Set %r8 to &p[0]
Set %rcx to &p[0] (%rcx = p)

```
long vframe(long n, long idx, long *q) {
    long i;
    long *p[n];
    p[0] = &i;
    for (i = 1; i < n; i++)
        p[i] = q;
    return *p[idx];
```

%rbp作为帧指针的作用：

- 在整个函数的执行过程中，%rbp始终指向函数栈的顶端（在返回地址和保存被调用者保存寄存器的值的下方）
- 利用固定长度的局部变量相对于%rbp的偏移量来引用它们
- leave指令释放整个栈帧

Set stack pointer to beginning of frame
Restore saved %rbp and set stack ptr
to end of caller's frame

Restore %rbp and %rsp
Return

如何防御堆栈溢出攻击

- 用fgets规定输入字符串的大小。gets很不安全！
- 给内存增加可执行权限标记，禁止执行栈上的代码
- 栈偏移量随机化，无法事先确定数据地址
- 堆栈金丝雀：`%fs:0x28`

(了解即可)

Memory Layout

■ Stack

- Runtime stack (8MB limit)
- E.g., local variables

■ Heap

- Dynamically allocated as needed
- When call `malloc()`, `calloc()`, `new()`

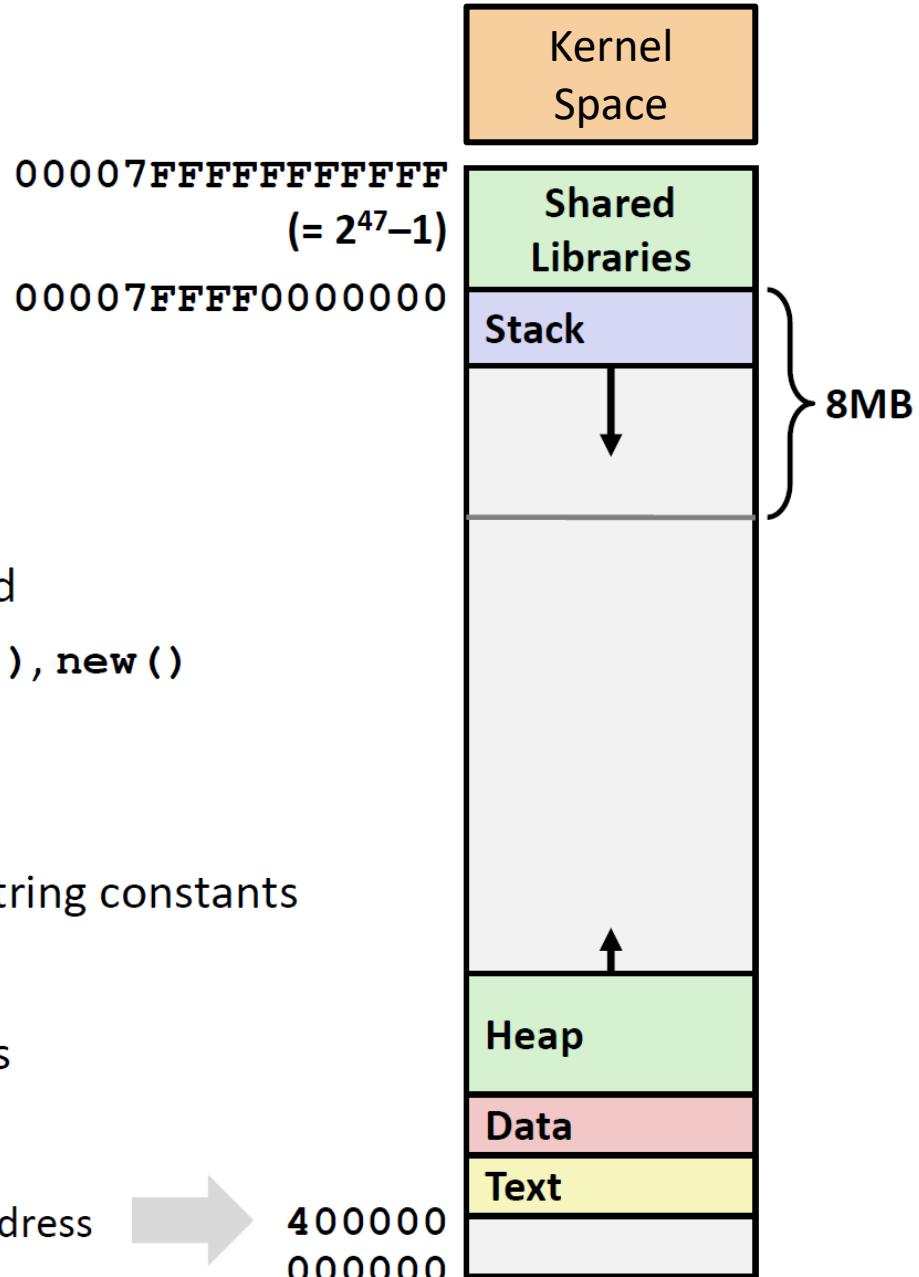
■ Data

- Statically allocated data
- E.g., global vars, `static` vars, string constants

■ Text / Shared Libraries

- Executable machine instructions
- Read-only

Hex Address → 400000
000000



Any Questions?
